

BRNANL, A Fortran Program to  
Identify Basic Blocks in Fortran Programs\*

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## Abstract

A basic block is a sequence of consecutive Fortran statements which must be executed consecutively; that is, if one statement in the block is executed, all are executed. Except for special cases noted in the text, a Fortran program is a catenation of basic blocks. BRNANL is a Fortran program designed to recognize basic blocks in a Fortran program. Given a Fortran program (FP) BRNANL will generate a modified Fortran program (MFP) in which a subroutine call is located at the head of every basic block. Execution of the MFP produces the same results as execution of the FP but the inserted subroutine calls permit monitoring of the execution sequences. User information for running BRNANL is presented.

Keywords: Software testing, control path analysis.



## 1. Introduction

This report describes external features of a program BRNANL which is designed to identify basic blocks in an ANSI Fortran [1] program. Informally, a basic block (BB) is a sequence of statements which must be executed consecutively: a precise definition will be given later. The following example serves to illustrate the idea of a BB in Fortran:

```
      .
      .
      .
      K = K + 1
      IF(K) 10, 20, 30 } tail end of BB
10 X = X + Y } BB
20 Y = 3.0*Z + 9.0
   D = A*B + C*D } BB
   GO TO(40, 50), J
40 A = 5.0 } head end of BB
      .
      .
      .
```

The notion of a basic block which we use is similar, but not identical, to that used in the specification of ANSI Fortran ([1], section 10.2.7). It follows more closely the definition generally used in the code optimization literature [2, 3, 4, 5]. Programs similar to this have appeared before. The program which is most similar to ours is one called FETE [6] written by Ingalls\*, another program of this type appears to have been contained in a larger program reported by Allen [4]. A slightly different, but related program has been reported by Russell and Estrin [7]. We desired a program of this type, written in ANSI Fortran for portability, which could be easily

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\* An improved version, called FORTUNE, is commercially available from Capex (Phoenix, Arizona).

modified to meet various needs we had in connection with a project on software validation. For these reasons we created the program described here.

BRNANL accepts as input a syntactically correct ANSI Fortran program, say FP, and produces a modified form of FP, say MFP, differing from FP in that a subroutine call has been placed at the beginning of every BB. BRNANL numbers the BBs in order of their appearance in the source code and this number appears as a parameter in the inserted subroutine call. There is a second calling parameter which is used to identify special situations. A subroutine call is inserted before the first executable statement of a program; the second parameter has the value 1 in this case. A subroutine call is inserted before every STOP statement of a program; the second parameter has the value 3 in this case. In the normal case, when the inserted call appears as the first statement of a BB the second parameter has the value 2. The MFP for the example above is shown below.

```
      .  
      .  
      .  
      K = K + 1  
      IF(K) 10, 20, 30  
10    CALL XXXXXX(5, 2)  
      X = X + Y  
20    CALL XXXXXX(6, 2)  
      Y = 3.0*Z + 9.0  
      D = A*B + C*D  
      GO TO (40, 50), J  
40    CALL XXXXXX(7, 2)  
      A = 5.0  
      .  
      .  
      .
```

Here the first BB of the segment has been arbitrarily numbered 5 and the called subroutine arbitrarily named XXXXXX.

The MFP produced by BRNANL has the physical form of a printed listing and/or a punched deck and/or a file which may be on disc or tape depending on the system under which it executes. All statements inserted by BRNANL in the MFP are flagged by asterisks in columns 73-80 of the output. The printed listing and the file have the block number of each statement and the sequential line number recorded in columns to the right of each statement. A copy of the FP and the MFP are shown in Appendix E for a subroutine subprogram.

It is possible to suppress entirely the insertion of the subroutine calls. This is controlled by a datum on a data card read by BRNANL (cf. Appendix A). When this option is used, the listing which is produced will have the block number and line number at the right of each statement as before. This option is used to analyze the structure of the flowgraph for the program. In this report we are primarily concerned with using BRNANL to obtain a MFP which does have the subroutine calls inserted, so no further consideration is given to this option.

When the MFP is executed, various types of information can be recorded depending on the subroutine XXXXXX\*. For example, the set of basic blocks executed can be recorded, the frequency of execution of basic blocks can be recorded, the sequence in which basic blocks are executed can be recorded, etc. Used in this way BRNANL is a valuable tool in program testing and it

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\* Henceforth we will use XXXXXX for the name of the subroutine appearing in the inserted call.

was with this purpose in mind that BRNANL was constructed. Since its construction, we have found it to be a useful tool in the reduction of a Fortran program to a directed graph.

Although BRNANL was written in ANSI Fortran it does contain one machine dependent subroutine CHRCHK which is designed to classify a character which has been read with an A1 format specification as a letter, a digit, or a special character. Specifications of this subroutine will be found in Appendix D.

Use of BRNANL is very simple. One card containing parameter specifications is placed in front of the FP and one card containing the character \$ in column 7 is placed in back of the FP: the resulting deck is the data deck for BRNANL. This deck setup is shown in Appendix B. The first card contains the following parameter specifications: name of the subroutine for the inserted call; a unique Fortran variable name (i.e. a name not used in the FP); initial block number; initial line number; flag to indicate suppression of inserted subroutine calls. Details are in Appendix A. The total storage required for the assembled program on the CDC 6400 computer operating under KRONOS 2.1 using the RUN compiler is 4240<sub>8</sub> words. In this environment sixty-six seconds of central processor time was required to run BRNANL on itself which consists of 1737 source statements excluding comments.



## 2. Basic Blocks

In this section the precise rules for identifying basic blocks and inserting subroutine calls are given. As already indicated, a BB is a sequence of one or more consecutive, executable statements in a source program. Suppose we identify the consecutive executable statements in a source program as  $S_1, S_2, \dots, S_n$ . Each BB consists of some subsequence, say  $S_j, S_{j+1}, \dots, S_{j+k}$ ; the first statement,  $S_j$  is called the head, the last statement,  $S_{j+k}$ , is called the tail, and the sequence of statements between the head and the tail,  $S_{j+1}, S_{j+2}, \dots, S_{j+k-1}$  is called the trunk. The trunk may be empty and the head and tail may be embodied in a single statement.

A tail is any one of the following:

- (a) logical IF statement;
- (b) arithmetic IF statement;
- (c) DO statement;
- (d) any form of GO TO statement;
- (e) RETURN statement
- (f) STOP statement;
- (g) any statement followed by a labelled statement except when the labelled statement is a FORMAT statement or when the labelled statement is the terminal statement in a DO loop;
- (h) the terminal statement in a DO loop.

A head is the statement immediately following a tail with two exceptions: the terminal statement in a DO-loop is never a head; the first executable statement of the main program and every subprogram is a head.

From this definition, excepting three special situations discussed below, the following assertions are true:

- (a) Every  $S_i$  belongs to a BB and cannot belong to more than one BB.
- (b) If any  $S_i$  in a BB is executed then every statement in that BB is executed.

From the first assertion it follows that we may view a program as a catenation of basic blocks. From the second assertion we may conclude that if the head statement of every BB of the program is executed then every statement of the program is executed.

### 3. Exceptional Situations

The logical IF statement presents one exceptional situation. Consider the statement:

```
IF(K.LT.0) X = X + Y.
```

This statement is a tail, however execution of this tail does not necessarily imply execution of the embedded assignment statement  $X = X + Y$ . To resolve this we treat only the structure

```
IF(<logical expression>)
```

as a tail and the portion of the logical IF following this is treated as a BB consisting of one statement. Thus in the above example IF(K.LT.0) would be the tail for, say, BB(12), then the statement  $X = X + Y$  would be BB(13). With this understanding the two assertions above remain valid.

The second exceptional situation arises when a jump within a DO-loop can go to the last statement in the scope of a DO, as illustrated in the following situation

```
DO 20 J = 1, K
X(J) = X(J) + Y
IF(X(J)) 10, 20, 30
10 L = L + 1
20 V(J) = 0
30 A = B + C
```

Since there is a jump possible to statement

```
20 V(J) = 0
```

we ought to identify it as a head -- it is a tail by virtue of it being the terminal statement in a DO-loop (rule (k) above). However, if we were to treat this statement also as a head, then we would have in the MFP

```
20 CALL XXXXXX(--,--)
V(J) = 0
```

violating the DO-loop. We resolve this problem by not permitting the terminal statement in a DO-loop to be a head; it is always treated only as a tail. The MFP for the program segment above is

```
DO 20 J = 1, K
  CALL XXXXXX(12, 2)
  X(J) = X(J) + Y
  IF(X(J)) 10, 20, 30
10  CALL XXXXXX(13, 2)
    L = L + 1
20  V(J) = 0
    CALL XXXXXX(14, 2)
30  CALL XXXXXX(15, 2)
    A = B + c
```

where the BB numbering arbitrarily starts at 12 and BB(14) is a dummy used to detect satisfying the DO-loop. Thus the pair of statements

```
10  L = L + 1
20  V(J) = 0
```

is BB(13) and it is evident that assertion (b) above is not true. On the other hand it is important to note that if the head of every BB (including the dummy) is executed, then it is true that every statement has been executed. Also the number of times the statement

```
20  V(J) = 0
```

is executed is given by the expression

$$n_{12} - n_{11} + n_{14}$$

where  $n_i$  is the number of times BB(i) is entered. Finally, we observe that if DO-loops are terminated with CONTINUE statements, a good programming practice in any case, then jumps to the end of a DO-loop do not cause any important difficulty since one is not usually interested in the execution of CONTINUE statements.

The third exceptional situation arises when there is no return to the calling program after a subprogram has been called into execution. For example, in the basic block,

```

    --
    --
    --
30   J = J + 1
      X = X + Y
      CALL XAMPL(X, J, Z)
      X = Z
      GO TO 20
    --
    --
    --

```

failure to return from XAMPL makes assertion (b) above false. If STOP statements are permitted only in the main program, then this situation cannot arise unless execution is aborted by the system due to an error (overflow, array bounds violation, etc.).

These special situations could be eliminated. The logical IF problem could be removed by a replacement of the logical IF by an arithmetic IF and suitable restructuring of the program. The DO-loop problem could be removed by permitting assignment of new statement labels and appropriate relabelling. Finally, a change in the rules defining a BB could partially eliminate the STOP statement problem. A subroutine CALL statement could be a BB but there would still be a problem with FUNCTION calls since these are embedded in statements.

The splitting of a logical IF statement in two BBs is done in the following way. Suppose we have the logical IF statement

<label> IF(<Boolean expression>)<statement>

then in the MFP this appears as

```
<label> LLLLLL = <Boolean expression>  
        IF(LLLLLL) CALL XXXXXX(--,--)  
        IF(LLLLLL) <statement>
```

It is evident that the subroutine call will be executed if and only if <statement> is executed so the call can be associated with the BB for the statement. The name LLLLLL is arbitrary: it is the second parameter on the data card. The MFP will also have a type declaration:

```
LOGICAL LLLLLL
```

Since a logical IF can terminate a DO-loop, it is evident that this situation presents a special problem. Splitting of the logical IF is not done in this case. Thus in the following sequence

```
        DO 10 I = 1, N  
        --  
        --  
        --  
10      IF(X.LT.0) X = 1
```

The "BB" X = 1 is not identified as a BB. If the assertion (b) is to be valid, it is evident that a logical IF terminating a DO-loop must be prohibited. When BRNANL detects this situation it prints an error message.

In our own use of BRNANL we preprocess the FP with another program, STYLE [8], which reformats the FP and causes each DO loop to terminate on a CONTINUE statement. This essentially removes the difficulties cited above.

4. First Executable Statement and STOP Statement.

The first BB executed in a program is given special treatment. Suppose the FP begins with the statements

```
C THIS IS THE MAIN PROGRAM
  DIMENSION A(10), B(10, 10)
10  READ (5, 999) A
    DO 20 I = 1, 10
      ---
      ---
      ---
```

Then the MFP begins with the statements

```
C THIS IS THE MAIN PROGRAM
  DIMENSION A(10),B(10, 10)
    CALL XXXXXX(0, 1)
10  CALL XXXXXX(1, 2)
    READ(5, 999) A
    DO 20 I = 1, 10
```

The second parameter of the first CALL is 1, uniquely identifying it as preceding the first executable statement in the program. This information allows the routine XXXXXX to perform initialization. The first parameter in this call is one less than the initial block number, the third parameter on the data card (cf. Appendix A); here it is assumed that this number was 1. The second CALL, which would be there even if the label were not on the READ statement, is used to identify actual entry into the BB. It is to be noted that this situation arises only in a main program.

The BB in which a STOP appears as the tail is treated in a special way. In addition to the call which is inserted at the head of the block, a call is inserted immediately before the STOP statement. The following example illustrates this. Suppose the FP contains the BB, say BB(100),

```
150  X = SIN(Y)
      WRITE(6, 999) X, Y
      STOP
```

Then the MFP is

```
150  CALL XXXXXX(100, 2)
      X = SIN(Y)
      WRITE(6, 999) X, Y
      CALL XXXXXX(100, 3)
      STOP
```

It is to be noted that the second parameter in the CALL just before the STOP is 3; this special value is used only before a STOP so the routine XXXXXX can take whatever steps are appropriate for such a condition.

Typically it would print accumulated data on BB activity in executing the MFP.



5. Limitations.

The most important limitation arises from the fact that BRNANL assumes that the FP is a syntactically correct ANSI Fortran program; if it is not, incorrect execution may result.

Other limitations are listed below:

1. Maximum number that can be assigned to a BB is 9999;
2. Maximum number of subscripted variables in each program unit (subroutine subprogram, function subprogram, main program) is 50;
3. Maximum depth for DO-loop nesting is 29.

It is recommended that all DO-loops terminate on CONTINUE statements. (Preprocessing the FP by STYLE [8] will guarantee this.) If a DO-loop does not terminate on a CONTINUE statement BRNANL will still execute properly, however jumps to the last statement in the DO need special consideration as described in section 3 of this report.

A list of error messages which can be produced by BRNANL is given in Appendix B.

6. I/O FILES.

BRNANL reads the input file from unit 5, writes the print file on unit 6, and writes the punch file on unit 7. Specifically all READ statements have the form

READ(KIN, ...

all WRITE statements for producing the listing of the MFP and any error messages have the form

WRITE(KPR, ...

and all write statements for producing the source "deck" for the MFP have the form

WRITE(KPU, ...

A DATA statement is used assign 5, 6, 7 to KIN, KPR, KPU, respectively.

7. Acknowledgements.

Much of the testing, corrections to errors in the original program, and insertion of COMMENTS was done by Jeffery Wright and Jacob Wu, graduate students in the Department of Computer Science at the University of Colorado.

8. References.

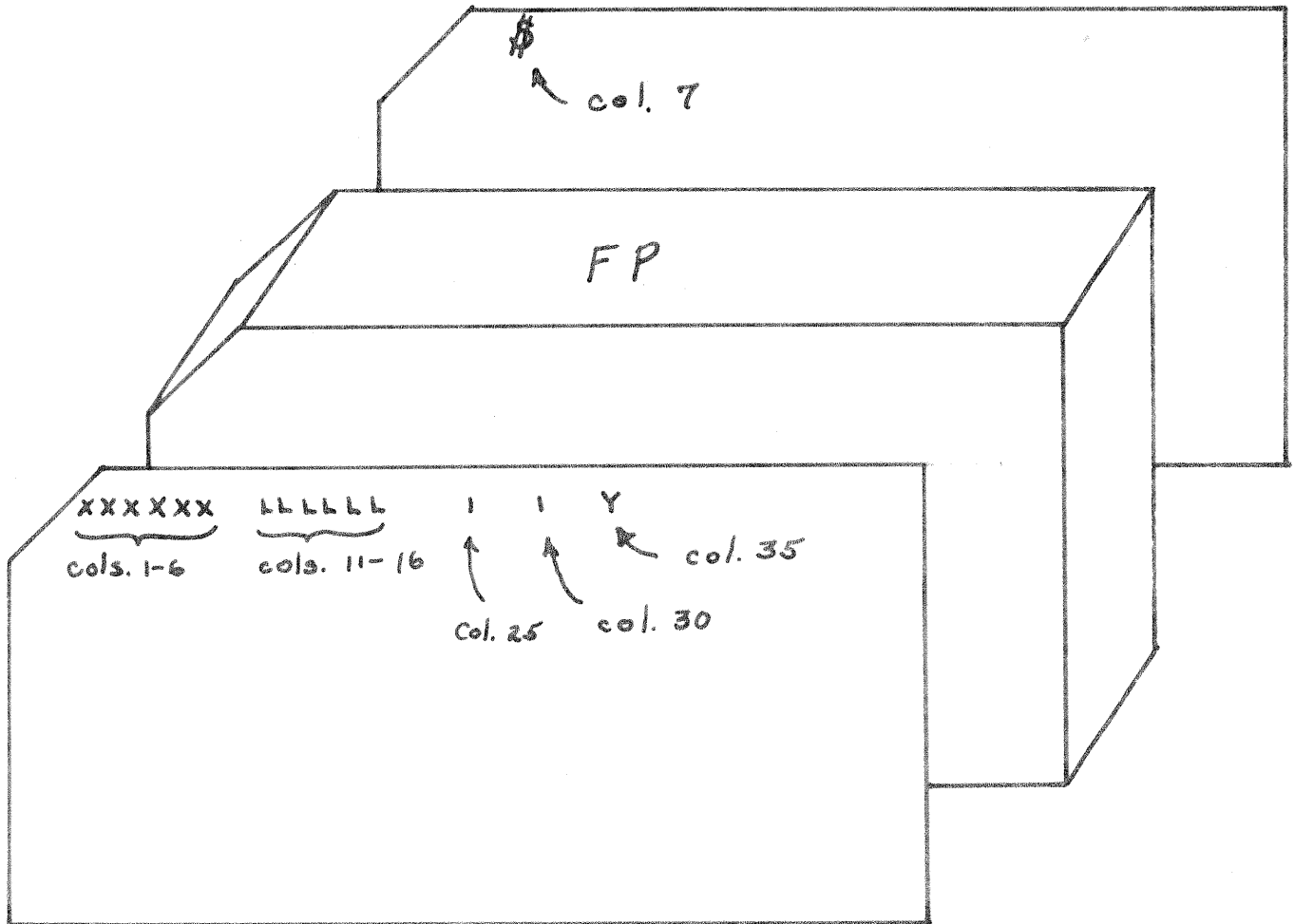
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See also:
  - (a) Fortran vs. Basic Fortran, Comm ACM 7 (Oct. 1964), 591-625.
  - (b) Clarification of Fortran Standards - Initial Progress, Comm A M 12 (May 1969), 289-294.
  - (c) Clarification of Fortran Standards - Second Report, Comm ACM 14 (Oct 1971), 628-642.
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7. E. C. Russell and G. Estrin, Measurement based automatic analysis of Fortran programs. SJCC (1969), 723-732.
8. Dorothy Lang Wedel, STYLE Editor: User's Guide, Department of Computer Science, University of Colorado, Report 7 (1972).

Appendix A: Data card.

One data card must precede the FP. The layout of this card follows:

<u>cols.</u>	<u>contents</u>
1-6	Name of subroutine appearing in the inserted CALL statements;
11-16	Name of variable used to hold value of Boolean expressions appearing in logical IF statements;
22-25	Number of first BB in program; BBs are numbered sequentially in order of appearance in the FP; value entered as a right-justified integer;
26-30	Number of first line for sequential numbering of lines in output file supplied by BRNANL; value entered as a right-justified integer.
35	Y if subroutine calls are to be inserted.

Appendix B: Run Deck Organization



Appendix C: Error Messages

1: MORE THAN 20 CARDS USED FOR STATEMENT.

The maximum number of cards permitted for a statement is 20 (i.e. 19 continuation cards). Fatal error.

2: IF STATEMENT SYNTAX ERROR.

ANSI Fortran syntax error. Fatal error.

3: MORE THAN 50 DIMENSIONED VARIABLES.

An array in BRNANL called ARRNAM holds the list of dimensioned variables in the subprogram, or main program being processed. It is dimensioned at 50. Fatal error.

4: STATEMENT ENDS WITH LEFT PARENTHESIS.

ANSI Fortran syntax error. Fatal error.

5: DECLARATION FOLLOWED ONLY BY BLANKS.

ANSI Fortran syntax error. Fatal error.

6: LIMIT OF 9999 ON BASIC BLOCK INDEX EXCEEDED.

BRNANL requires BB index to lie in the interval (1, 9999). Fatal error.

7: NON-ANSI BLANK CARD ENCOUNTERED.

Blank cards are not permitted in the subject program. This card is automatically replaced by a blank comment card. Non-fatal error.

8: NON-ANSI PROGRAM CARD ENCOUNTERED.

A PROGRAM card is required for the main program in CDC Fortran, however this is not legal ANSI Fortran. This card is ignored. Non-fatal error.

9: LABEL CONTAINS AN ILLEGAL CHARACTER.

A label must consist of digits only. Fatal error.

10: LOGICAL IF CLOSING A DO-LOOP.

After each error message the following information is printed:

BUFFER A CONTAINS

---(statement being processed)

BUFFER B CONTAINS

---(next statement to be processed)

BUFFER D CONTAINS

---(first card of next statement)



Appendix D: Machine Dependent Subroutine

The subroutine CHRCHK in BRNANL is machine dependent and may have to be modified by the user for systems other than the CDC 6400. This subroutine determines whether a character is a letter, a digit, or special.

The subroutine specification is

```
SUBROUTINE CHRCHK(A, I, L)
```

where the formal parameters are defined as follows:

A -- a one dimensional array holding characters which have been read into it using an A1 format specification.

I -- the character to be checked is in position A(I).

L -- the routine CHRCHK makes the assignment

L = 1 if the character in A(I) is a letter

L = 2 if the character in A(I) is a digit

L = 3 if the character in A(I) is special.

A listing of this subroutine is on the following page.

```
      SUBROUTINE CHRCHK(A, I, L)
C THIS IS A CHARACTER CHECK ROUTINE FOR BUFFER A. L=1 IF
C A(I) IS A LETTER, L=2 IF A(I) IS A DIGIT,L=3 IF A(I) IS A
C SPECIAL CHARACTER.
      DIMENSION A(I)
      INTEGER ALPBL, ALPBH, NUML, NUMH, A
C THE FOLLOWING CHECKING TECHNIQUE SHOULD BE ADJUSTED FOR
C LOCAL CHARACTER SET IF ALPHABET CHARACTERS OR NUMERIC
C CHARACTERS ARE NOT
C IN A CONTIGUOUS GROUP.
      DATA ALPBL /1HA/, ALPBH /1HZ/, NUML /1H0/, NUMH /1H9/
      ICHR = A(I)
      IF (ICHR-ALPBL) 20, 10, 10
10  IF (ICHR-ALPBH) 40, 40, 20
20  IF (ICHR-NUML) 60, 30, 30
30  IF (ICHR-NUMH) 50, 50, 60
40  L = 1
      RETURN
50  L = 2
      RETURN
60  L = 3
      RETURN
      END
```

Appendix E: Example

On the following six pages an example illustrating the output obtained from BRNANL is shown. The first three pages contain the listing of the FP (a subroutine subprogram KZEOE). The next three pages contain the MFP as contained on the print file. In this example block numbering starts at 1 and line numbering start, at 10. Block numbers associated with logical IF statements are flagged by an asterisk.

```

SUBROUTINE KZEONE(X, Y, RE0, IM0, RE1, IM1)
C THE VARIABLES X AND Y ARE THE REAL AND IMAGINARY PARTS OF
C THE ARGUMENT OF THE FIRST TWO MODIFIED BESSEL FUNCTIONS
C OF THE SECOND KIND, K0 AND K1. RE0, IM0, RE1 AND IM1 GIVE
C THE REAL AND IMAGINARY PARTS OF EXP(X)*K0 AND EXP(X)*K1,
C RESPECTIVELY. ALTHOUGH THE REAL NOTATION USED IN THIS
C SUBROUTINE MAY SEEM INELEGANT WHEN COMPARED WITH THE
C COMPLEX NOTATION THAT FORTRAN ALLOWS, THIS VERSION RUNS
C ABOUT 30 PERCENT FASTER THAN ONE WRITTEN USING COMPLEX
C VARIABLES.
      DOUBLE PRECISION X, Y, X2, Y2, RE0, IM0, RE1, IM1,
      * R1, R2, T1, T2, P1, P2, RTERM, ITERM, EXSQ(8), TSQ(8)
      DATA TSQ(1) /0.000/, TSQ(2) /3.143036339206350-1/,
      * TSQ(3) /4.2907586229591500/, TSQ(4)
      * /2.9583744586966500/, TSQ(5) /5.40990315972444400/,
      * TSQ(6) /8.8040795780567600/, TSQ(7)
      * /1.3468535743251501/, TSQ(8) /2.0249916365870901/,
      * EXSQ(1) /0.564100308726400/, EXSQ(2)
      * /0.412028687498900/, EXSQ(3) /0.158488915795900/,
      * EXSQ(4) /0.30780033872550-1/, EXSQ(5)
      * /0.2778008429130-2/, EXSQ(6) /0.10000444123250-3/,
      * EXSQ(7) /0.10591155477110-5/, EXSQ(8)
      * /0.15224758042540-d/
C THE ARRAYS TSQ AND EXSQ CONTAIN THE SQUARE OF THE
C ABSCESSAS AND THE WEIGHT FACTORS USED IN THE GAUSS-
C HERMITE QUADRATURE.
      P2 = X*X + Y*Y
      IF (X.GT.0.000 .OR. R2.NE.0.000) GO TO 10
      WRITE (6,'99999')
      RETURN
      10 IF (R2.GE.1.9602) GO TO 50
      IF (R2.GE.1.84901) GO TO 30
C THIS SECTION CALCULATES THE FUNCTIONS USING THE SERIES
C EXPANSIONS
      X2 = X/2.000
      Y2 = Y/2.000
      P1 = X2*X2
      P2 = Y2*Y2
      T1 = -(DLG(P1+P2)/2.000+0.577215664901532900)
C THE CONSTANT IN THE PRECEDING STATEMENT IS EULER*S
C CONSTANT
      T2 = -DATAN2(Y,X)
      X2 = P1 - P2
      Y2 = X*Y2
      RTERM = 1.000
      ITERM = 0.000
      RE0 = T1
      IM0 = T2
      T1 = T1 + 0.500
      RE1 = T1
      IM1 = T2
      P2 = DSQRT(R2)
      L = 2.10600*P2 + 4.400
      IF (P2.LT.8.00-1) L = 2.12900*P2 + 4.000
      DO 20 N=1,L
      P1 = N
      P2 = N*N
      R1 = RTERM
      RTERM = (R1*X2-ITERM*Y2)/P2
      ITERM = (R1*Y2+ITERM*X2)/P2
      T1 = T1 + 0.500/P1
      RE0 = RE0 + T1*ITERM
      IM0 = IM0 + T1*ITERM + T2*ITERM
      KZE 10
      KZE 20
      KZE 30
      KZE 40
      KZE 50
      KZE 60
      KZE 70
      KZE 80
      KZE 90
      KZE 100
      KZE 110
      KZE 120
      KZE 130
      KZE 140
      KZE 150
      KZE 160
      KZE 170
      KZE 180
      KZE 190
      KZE 200
      KZE 210
      KZE 220
      KZE 230
      KZE 240
      KZE 250
      KZE 260
      KZE 270
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      KZE 290
      KZE 300
      KZE 310
      KZE 320
      KZE 330
      KZE 340
      KZE 350
      KZE 360
      KZE 370
      KZE 380
      KZE 390
      KZE 400
      KZE 410
      KZE 420
      KZE 430
      KZE 440
      KZE 450
      KZE 460
      KZE 470
      KZE 480
      KZE 490
      KZE 500
      KZE 510
      KZE 520
      KZE 530
      KZE 540
      KZE 550
      KZE 560
      KZE 570
      KZE 580
      KZE 590
      KZE 600
      KZE 610
      KZE 620
      KZE 630

```

KZE 640  
 KZE 650  
 KZE 660  
 KZE 670  
 KZE 680  
 KZE 690  
 KZE 700  
 KZE 710  
 KZE 720  
 KZE 730  
 KZE 740  
 KZE 750  
 KZE 760  
 KZE 770  
 KZE 780  
 KZE 790  
 KZE 800  
 KZE 810  
 KZE 820  
 KZE 830  
 KZE 840  
 KZE 850  
 KZE 860  
 KZE 870  
 KZE 880  
 KZE 890  
 KZE 900  
 KZE 910  
 KZE 920  
 KZE 930  
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 KZE1170  
 KZE1180  
 KZE1190  
 KZE1200  
 KZE1210  
 KZE1220  
 KZE1230  
 KZE1240  
 KZE1250  
 KZE1260

```

PI = PI + 1.000
TI = TI + 0.500/PI
REI = REI + (T1*ITERM - T2*ITERM)/PI
IMI = IMI + (T1*ITERM + T2*ITERM)/PI
20 CONTINUE
R1 = X/R2 - 0.500*(A*REI - Y*IMI)
R2 = -Y/R2 - 0.500*(X*IMI + Y*REI)
PI = DEXP(X)
REO = PI*REO
IM0 = PI*IM0
REI = PI*RI
IMI = PI*R2
RETURN
C THIS SECTION CALCULATES THE FUNCTIONS USING THE INTEGRAL
C REPRESENTATION, EGN 3, EVALUATED WITH 15 POINT GAUSS-
C HERMITE QUADRATURE
30 X2 = 2.000*X
Y2 = 2.000*Y
R1 = Y2*Y2
P1 = USQRT(X2*X2+R1)
P2 = USQRT(PI*X2)
T1 = EXSQ(1)/(2.000*PI)
REO = T1*P2
IM0 = T1/P2
REI = 0.000
IMI = 0.000
DO 40 N=2,8
T2 = X2 + TSQ(N)
P1 = USQRT(T2*T2+R1)
P2 = USQRT(PI+T2)
T1 = EXSQ(N)/PI
REO = REO + T1*P2
IM0 = IM0 + T1/P2
T1 = EXSQ(N)*TSQ(N)
REI = REI + T1*P2
IMI = IMI + T1/P2
40 CONTINUE
T2 = -Y2*IM0
REI = REI/R2
R2 = Y2*IM1/R2
RTERM = 1.4142135623730900*DCOS(Y)
ITERM = -1.4142135623730900*DSIN(Y)
C THE CONSTANT IN THE PREVIOUS STATEMENTS IS*OF COURSE*
C SORT(2.0).
IM0 = REO*ITERM + T2*ITERM
REO = REO*ITERM - T2*ITERM
T1 = REI*ITERM - R2*ITERM
T2 = REI*ITERM + R2*ITERM
REI = T1*A + T2*Y
IMI = -T1*Y + T2*X
RETURN
C THIS SECTION CALCULATES THE FUNCTIONS USING THE
C ASYMPTOTIC EXPANSIONS
50 RTERM = 1.000
ITERM = 0.000
REO = 1.000
IM0 = 0.000
REI = 1.000
IMI = 0.000
P1 = 8.000*R2
P2 = USQRT(R2)
L = 3.9100+8.1201/P2
R1 = 1.000

```

KZE1270  
 KZE1280  
 KZE1290  
 KZE1300  
 KZE1310  
 KZE1320  
 KZE1330  
 KZE1340  
 KZE1350  
 KZE1360  
 KZE1370  
 KZE1380  
 KZE1390  
 KZE1400  
 KZE1410  
 KZE1420  
 KZE1430  
 KZE1440  
 KZE1450  
 KZE1460  
 KZE1470  
 KZE1480  
 KZE1490  
 KZE1500  
 KZE1510  
 KZE1520  
 KZE1530  
 KZE1540  
 KZE1550  
 KZE1560  
 KZE1570  
 KZE1580  
 KZE1590  
 KZE1600  
 KZE1610

```

R2 = 1.0D0
M = 8
K = 3
CO 60 N=1*L
M = M + 8
K = K - M
R1 = FLOAT(K-4)*PI
R2 = FLOAT(K)*R2
T1 = FLOAT(N)*PI
T2 = RTERM
RTERM = (T2*X+ITERM*Y)/T1
ITERM = (-T2*Y+ITERM*X)/T1
R0 = R0 + R1*RTERM
IM0 = IM0 + R1*ITERM
RE1 = RE1 + R2*RTERM
IMI = IM1 + R2*ITERM
60 CONTINUE
T1 = DSQRT(P2*X)
T2 = -Y/T1
PI = 8.86226925427580-1/P2
C THIS CONSTANT IS SQRT(PI)/2.0, WITH PI=3.14159...
RTERM = PI*DCOS(Y)
ITERM = -PI*DSIN(Y)
R1 = R0*RTERM - IM0*ITERM
R2 = R0*ITERM + IM0*RTERM
RE0 = T1*R1 - T2*R2
IM0 = T1*R2 + T2*R1
R1 = RE1*RTERM - IM1*ITERM
R2 = RE1*ITERM + IM1*RTERM
RE1 = T1*R1 - T2*R2
IM1 = T1*R2 + T2*R1
RETURN
FORMAT (42H ARGUMENT OF THE BESSEL FUNCTIONS IS ZERO,
* 35H OR LIES IN LEFT HALF COMPLEX PLANE)
END
  
```

```

SUBROUTINE KZEONE(X, Y, RE0, IM0, RE1, IM1)
LOGICAL LLLLLL
C THE VARIABLES X AND Y ARE THE REAL AND IMAGINARY PARTS OF
C THE ARGUMENT OF THE FIRST TWO MODIFIED BESSEL FUNCTIONS
C OF THE SECOND KIND, K0 AND K1. RE0, IM0, RE1 AND IM1 GIVE
C THE REAL AND IMAGINARY PARTS OF EXP(X)*K0 AND EXP(X)*K1,
C RESPECTIVELY. ALTHOUGH THE REAL NOTATION USED IN THIS
C SUBROUTINE MAY SEEM INELEGANT WHEN COMPARED WITH THE
C COMPLEX NOTATION THAT FORTRAN ALLOWS, THIS VERSION RUNS
C ABOUT 30 PERCENT FASTER THAN ONE WRITTEN USING COMPLEX
C VARIABLES.
      DOUBLE PRECISION X, Y, X2, Y2, RE0, IM0, RE1, IM1,
      * P1, R2, T1, T2, P1, P2, RTERM, ITERM, EXSQ(8), TSQ(8)
      DATA TSQ(1) /0.000/, TSQ(2) /3.193036339206350-1/,
      * TSQ(3) /1.290758629591500/, TSQ(4)
      * /2.9583744586966500/, TSQ(5) /5.4090315972444400/,
      * TSQ(6) /8.8040795780567600/, TSQ(7)
      * /1.3468535743251501/, TSQ(8) /2.0249916365870901/,
      * EXSQ(1) /0.564100308/26400/, EXSQ(2)
      * /0.412028687498900/, EXSQ(3) /0.158488915795900/,
      * EXSQ(4) /0.307800338/2550-1/, EXSQ(5)
      * /0.27780688+29130-2/, EXSQ(6) /0.10000+44123250-3/,
      * EXSQ(7) /0.10591155+77110-5/, EXSQ(8)
      * /0.152247580425+0-8/
C THE ARRAYS TSQ AND EXSQ CONTAIN THE SQUARE OF THE
C ABSCISSAS AND THE WEIGHT FACTORS USED IN THE GAUSS-
C HERMITE QUADRATURE.
      CALL XXXXXX(10001,2)
      R2 = X*X + Y*Y
      LLLLLL=X.GT.0.000.UK.R2.NE.0.000
      IF(LLLLLL)CALL XXXXXX(0002,2)
      GO TO 10
      IF(LLLLLL)
      CALL XXXXXX(10003,2)
      *RITE (6,999999)
      RETURN
10 CALL XXXXXX(10004,2)
      LLLLLL=R2.GE.1.9602
      IF(LLLLLL)CALL XXXXXX(0005,2)
      IF(LLLLLL)
      GO TO 50
      CALL XXXXXX(10006,2)
      LLLLLL=R2.GE.1.84901
      IF(LLLLLL)CALL XXXXXX(0007,2)
      IF(LLLLLL)
      GO TO 30
C THIS SECTION CALCULATES THE FUNCTIONS USING THE SERIES
C EXPANSIONS
      CALL XXXXXX(10008,2)
      X2 = X/2.000
      Y2 = Y/2.000
      P1 = X2*X2
      P2 = Y2*Y2
      T1 = -(DL05(P1+P2)/2.000+0.5772156649015329D0)
C THE CONSTANT IN THE PRECEDING STATEMENT IS EULERS
C CONSTANT
      T2 = -DATAN2(Y,X)
      X2 = P1 - P2
      Y2 = X*Y2
      RTERM = 1.000
      ITERM = 0.000
      RE0 = Y1
      IM0 = T2
      T1 = T1 + 0.500

```

KZE 10	0	0
***	0	0
KZE 20	0	0
KZE 30	0	0
KZE 40	0	0
KZE 50	0	0
KZE 60	0	0
KZE 70	0	0
KZE 80	0	0
KZE 90	0	0
KZE 100	0	0
KZE 110	0	0
KZE 120	0	0
KZE 130	0	0
KZE 140	0	0
KZE 150	0	0
KZE 160	0	0
KZE 170	0	0
KZE 180	0	0
KZE 190	0	0
KZE 200	0	0
KZE 210	0	0
KZE 220	0	0
KZE 230	0	0
KZE 240	0	0
KZE 250	0	0
KZE 260	0	0
*****	1	1
KZE 270	1	1
*****	1	1
*****	2	2
KZE 280	2	2
*****	3	3
KZE 290	3	3
KZE 300	3	3
*****	4	4
*****	4	4
*****	5	5
KZE 310	5	5
*****	6	6
*****	6	6
*****	7	7
KZE 320	7	7
KZE 330	7	7
KZE 340	7	7
*****	8	8
KZE 350	8	8
KZE 360	8	8
KZE 370	8	8
KZE 380	8	8
KZE 390	8	8
KZE 400	8	8
KZE 410	8	8
KZE 420	8	8
KZE 430	8	8
KZE 440	8	8
KZE 450	8	8
KZE 460	8	8
KZE 470	8	8
KZE 480	8	8
KZE 490	8	8

```

REL = T1
IMI = T2
P2 = DSQRT(R2)
L = 2.10600*P2 + 4.400
LLLL=P2*LT.8.00-1
IF(LLLL)CALL XXXXX(0009,2)
CALL XXXXX(0010,2)
DO 20 N=1,L
CALL XXXXX(0011,2)
PI = N
P2 = N*N
R1 = RTERM
RTERM = (R1*X2-ITERM*Y2)/P2
ITERM = (R1*Y2+ITERM*X2)/P2
T1 = T1 + 0.500/P1
RE0 = RE0 + T1*RTERM - T2*ITERM
IM0 = IM0 + T1*ITERM + T2*RTERM
PI = PI + 1.000
T1 = T1 + 0.500/P1
REL = REL + (T1*RTERM-T2*ITERM)/PI
IMI = IMI + (T1*ITERM+T2*RTERM)/PI
20 CONTINUE
CALL XXXXX(0012,2)
R1 = X/R2 - 0.500*(X*REL-Y*IMI)
R2 = -Y/R2 - 0.500*(X*IMI+Y*REL)
PI = DEAP(X)
RE0 = PI*RE0
IM0 = PI*IM0
REL = PI*R1
IMI = PI*R2
RETURN
C THIS SECTION CALCULATES THE FUNCTIONS USING THE INTEGRAL
C REPRESENTATION, EON 3, EVALUATED WITH 15 POINT GAUSS-
C HERMITE QUADRATURE
30 CALL XXXXX(0013,2)
X2 = 2.000*X
Y2 = 2.000*Y
R1 = Y2*Y2
P1 = DSQRT(X2*X2+R1)
P2 = DSQRT(P1*X2)
T1 = EXSQ(1)/(2.000*P1)
RE0 = T1*P2
IM0 = T1/P2
REL = 0.000
IMI = 0.000
DO 40 N=2,8
CALL XXXXX(0014,2)
T2 = X2 + ISQ(N)
P1 = DSQRT(T2*T2+R1)
P2 = DSQRT(P1*T2)
T1 = EXSQ(N)/P1
RE0 = RE0 + T1*P2
IM0 = IM0 + T1/P2
T1 = EXSQ(N)*ISQ(N)
REL = REL + T1*P2
IMI = IMI + T1/P2
40 CONTINUE
CALL XXXXX(0015,2)
T2 = -Y2*IM0
REL = REL/R2
R2 = Y2*IMI/R2
RTERM = 1.4142135623730900*DCOS(Y)

```

```

KZE 500 8 71
KZE 510 8 72
KZE 520 8 73
KZE 530 8 74
***** 8 75
***** 9* 76
KZE 540 9* 77
***** 10 78
KZE 550 10 79
***** 11 80
KZE 560 11 81
KZE 570 11 82
KZE 580 11 83
KZE 590 11 84
KZE 600 11 85
KZE 610 11 86
KZE 620 11 87
KZE 630 11 88
KZE 640 11 89
KZE 650 11 90
KZE 660 11 91
KZE 670 11 92
KZE 680 11 93
***** 12 94
KZE 690 12 95
KZE 700 12 96
KZE 710 12 97
KZE 720 12 98
KZE 730 12 99
KZE 740 12 100
KZE 750 12 101
KZE 760 12 102
KZE 770 12 103
KZE 780 12 104
KZE 790 12 105
***** 13 106
KZE 800 13 107
KZE 810 13 108
KZE 820 13 109
KZE 830 13 110
KZE 840 13 111
KZE 850 13 112
KZE 860 13 113
KZE 870 13 114
KZE 880 13 115
KZE 890 13 116
KZE 900 13 117
***** 14 118
KZE 910 14 119
KZE 920 14 120
KZE 930 14 121
KZE 940 14 122
KZE 950 14 123
KZE 960 14 124
KZE 970 14 125
KZE 980 14 126
KZE 990 14 127
KZE1000 14 128
***** 15 129
KZE1010 15 130
KZE1020 15 131
KZE1030 15 132
KZE1040 15 133

```



```

KZE1090 15
KZE1060 15
KZE1070 15
KZE1080 15
KZE1090 15
KZE1100 15
KZE1110 15
KZE1120 15
KZE1130 15
KZE1140 15
KZE1150 15
KZE1160 15
*****
KZE1170 16
KZE1180 16
KZE1190 16
KZE1200 16
KZE1210 16
KZE1220 16
KZE1230 16
KZE1240 16
KZE1250 16
KZE1260 16
KZE1270 16
KZE1280 16
KZE1290 16
KZE1300 16
*****
KZE1310 17
KZE1320 17
KZE1330 17
KZE1340 17
KZE1350 17
KZE1360 17
KZE1370 17
KZE1380 17
KZE1390 17
KZE1400 17
KZE1410 17
KZE1420 17
KZE1430 17
*****
KZE1440 18
KZE1450 18
KZE1460 18
KZE1470 18
KZE1480 18
KZE1490 18
KZE1500 18
KZE1510 18
KZE1520 18
KZE1530 18
KZE1540 18
KZE1550 18
KZE1560 18
KZE1570 18
KZE1580 18
KZE1590 18
KZE1600 18
KZE1610 18

```

```

ITEM = -1.4142135623730900*DSIN(Y)
C THE CONSTANT IN THE PREVIOUS STATEMENTS IS OF COURSE*
C SORT(2.0).
IM0 = REU*ITERM + I2*ITERM
REU = REU*ITERM - I2*ITERM
I1 = REI*ITERM - R2*ITERM
I2 = REI*ITERM + R2*ITERM
REI = I1*X + I2*Y
IM1 = -I1*Y + I2*X
RETURN

```

C THIS SECTION CALCULATES THE FUNCTIONS USING THE

```

C ASYMPTOTIC EXPANSIONS
50 CALL XXXXX(0016,2)
KTERM = 1.000
ITERM = 0.000
RE0 = 1.000
IM0 = 0.000
RE1 = 1.000
IM1 = 0.000
P1 = 6.000*R2
P2 = DSQRT(R2)
L = 3.9100+8.1201/P2
R1 = 1.000
R2 = 1.000
M = -8
K = 3
00 60 N=1,L
CALL XXXXX(0017,2)
M = M + 1
K = K - M
R1 = FLOAT(K-4)*R1
R2 = FLOAT(K)*R2
I1 = FLOAT(N)*P1
I2 = ITERM
ITERM = (I2*X+ITERM*Y)/I1
RE0 = RE0 + R1*ITERM
IM0 = IM0 + R1*ITERM
RE1 = RE1 + R2*ITERM
IM1 = IM1 + R2*ITERM

```

60 CONTINUE

```

CALL XXXXX(0018,2)
I1 = DSQRT(P2*X)
I2 = -Y/I1
PI = 8.862269254527580-1/P2
C THIS CONSTANT IS SQRT(PI)/2.0, WITH PI=3.14159...
ITERM = PI*DCOS(Y)
ITERM = -PI*DSIN(Y)
R1 = REU*ITERM - IM0*ITERM
R2 = REU*ITERM + IM0*ITERM
RE0 = I1*R1 - I2*R2
IM0 = I1*R2 + I2*R1
R1 = REI*ITERM - IM1*ITERM
R2 = REI*ITERM + IM1*ITERM
RE1 = I1*R1 - I2*R2
IM1 = I1*R2 + I2*R1
RETURN

```

99999 FORMAT (42H ARGUMENT OF THE BESSEL FUNCTIONS IS ZERO,  
\* 35H OR LIES IN LEFT HALF COMPLEX PLANE)  
END